

## MATHEMATICAL LOGIC — ASSIGNMENT TWO

(1) Prove  $\vdash (\forall x. A) \vee B \supset \forall x. A \vee B$  where  $x \notin \text{FV}(B)$ . Show a counterexample when  $x \in \text{FV}(B)$ .

$$\begin{array}{c}
 \dfrac{[\forall x. A]^2}{A} \forall E \\
 \dfrac{A \quad [B]^2}{A \vee B} \vee I_1 \quad \dfrac{[B]^2}{A \vee B} \vee I_2 \\
 \dfrac{A \vee B}{\forall x. A \vee B} \forall I \\
 \dfrac{\forall x. A \vee B}{(\forall x. A) \vee B \supset \forall x. A \vee B} \supset I^1
 \end{array}$$

Let  $A \equiv \perp$ ,  $B \equiv 'x \text{ is even}'$  in arithmetic, and suppose to take an evaluation of variables in which  $x \mapsto 2$ . Then  $\llbracket \forall x. A \vee B \rrbracket = \perp$  while  $\llbracket (\forall x. A) \vee B \rrbracket = \top$ .

(2) State and prove the Compactness Theorem.

This is Theorem 12.1 in the slides.

(3) In the definition of  $\Sigma$ -structure it is required that all sorts are interpreted in non-empty sets. Show that this requirement is necessary since the Soundness Theorem becomes invalid if the requirement is not met.

Consider a single-sorted signature with equality, and a  $\Sigma$ -structure  $\mathfrak{E}$  interpreting the sort in the empty set. By the following proof

$$\begin{array}{c}
 \dfrac{}{\forall x. x = x} \text{refl} \\
 \dfrac{\forall x. x = x}{x = x} \forall E \\
 \dfrac{x = x}{\exists x. x = x} \exists I
 \end{array}$$

assuming the Soundness Theorem to hold, it follows that  $\llbracket \exists x. x = x \rrbracket = 1$  in every model. So, in every model there is an element  $\llbracket x \rrbracket$  such that  $\llbracket x = x \rrbracket = 1$ . But this is false in  $\mathfrak{E}$ , since there is no element in the interpretation of the sort.

Hence,  $\exists x. x = x$  is provable but not valid in every model, contradicting the validity of the Soundness Theorem.